

# INFORMATION THEORY IN CS

Note Title

1/25/2016

## Objectives of this course

- Teach some basic Information Theory;
- Illustrate (mathematical) power by applying it in combinatorics, complexity theory, algorithms, privacy, ...

———— x ————

- Such diversity? How can one person know it all & teach it? ... Err... I can't + won't.
- Course will be run seminar style. I will cover first few lectures, rest will be student run. We will learn together.
- Grades based on?
  - Projects
  - Presentations
  - Participation
  - Scribe work

## Motivating Example: Shearer's Lemma

- Roughly relates volume of 3-d object to areas of its 2-d projections.
- Formally: sets in finite universe;  
cardinalities & not volume/area.

### Notation

- $[n] = \{1, 2, 3, \dots, n\}$
- for  $S = \{i_1, i_2, \dots, i_k\} \subseteq [d]$   
&  $x = (x_1, x_2, \dots, x_d) \in [n]^d$   
 $x_S \triangleq (x_{i_1}, x_{i_2}, \dots, x_{i_k})$
- for  $F \subseteq [n]^d$  &  $S \subseteq [d]$   
 $F_S \triangleq \{x_S \mid x \in F\}$

# Shearer's Lemma $(k, d)$ -version

Let  $F \subseteq [n]^d$ , then

$$|F| \binom{d-1}{k-1} \leq \left( \prod_{\substack{S \subseteq [d] \\ |S|=k}} |F_S| \right)$$

—————  $\times$  —————

- $(k, 1)$  version:

$$|F| \leq |F_{\{1,3\}}| \cdot |F_2| \cdot |F_3| \cdots |F_k|$$

Proof: trivial

$$F \subseteq F_1 \times F_2 \times F_3 \times \cdots \times F_k$$

- $(3, 2)$  version

$$|F|^2 \leq |F_{12}| \cdot |F_{23}| \cdot |F_{31}|$$

Already non-trivial!

# Proof via Information Theory

- Setup: - You & I know  $F$ .
  - I am given  $w \in F$
  - I need to describe  $w$  to you
  - How many bits do I need to send to you?

- Roughly  $H(w)$  measures this quantity

entropy of  $w$  ← key player!

- Notes: Entropy depends not only on  $F$ , but also how  $w$  is chosen from  $F$ .

• Suppose further  $W = (x, y, z)$  & you and I already know  $z$ , but Eve doesn't know  $z$ .

• How many bits does Eve expect me to send you?

•  $H(W|Z)$  denotes this quantity



Conditional Entropy of  $w$  given  $z$



key player #2

Aside: Today everything { definitions, claims, proofs } will be handwavy.

We will formalize everything later.

But today we will see why such formalism may be useful.

# Axioms Of (Conditional) Entropy

- (what we should expect; and will later confirm)

- $w \in F \Rightarrow H(w) \leq \log_2 |F|$

- $w$  uniform from  $F$   $\Rightarrow |F| = 2^j$

- $\Rightarrow H(w) = j = \log_2 |F|$

will ignore

- Suppose  $w = (x, y, z)$

$$H(x|y) \leq H(x)$$

(in the worst case, you and I can ignore

$y$  & focus on sending  $x$ .)

So  $x$  given  $y$  is easier to communicate.)

- $w = (x, y, z)$

$$H(x, y) = H(x) + H(y|x)$$

$$= H(y) + H(x|y)$$

Remarkable fact:

- $H(x)$  satisfying above axioms exists!
- Will see in lectures 2, 3 ...

————— x —————

Return to Shearer

(2,1) - case:

$$|F| \leq |F_1| \cdot |F_2|$$

$$\log |F| \leq \log |F_1| + \log |F_2|$$

Pick  $(x,y)$  uniformly from  $F$

- $H(x,y) = \log |F|$
- $\log |F_1| \geq H(x)$  (not equal since projecting  $(x,y) \rightarrow x$  is not uniform on  $F_1$ )
- $\log |F_2| \geq H(y)$

Suffices to show

$$H(X, Y) \leq H(X) + H(Y)$$

$$H(X, Y) = H(X) + H(Y|X)$$

$$\leq H(X) + H(Y) \quad \square$$

~~\_\_\_\_\_~~

(3,2) version of Shearer

wish to prove

$$|F|^2 \leq |F_{12}| \cdot |F_{23}| \cdot |F_{31}|$$

$\Uparrow$

$$2H(X, Y, Z) \leq H(X, Y) + H(Y, Z) + H(X, Z)$$

$$H(X, Y) = H(X) + H(Y|X)$$

RHS

$$H(Y, Z) = H(Y) + H(Z|Y)$$

$$H(X, Z) = H(X) + H(Z|X)$$

LHS

$$2H(X, Y, Z) = 2H(X) + 2H(Y|X) + 2H(Z|X, Y)$$

**QED**



# Proof of $(k, d)$ -Shearer

## Exercise 😊

- (Roughly: - Order variables  
- always condition on smaller # of variables  
- Condition on all smaller # of variables to get LHS)

## Moral of the story

- Somewhere in the axioms lies a very interesting inequality.
- Very<sup>n</sup> useful!
- IT has many other notions (Information, Divergence, Hellinger distance)
- Many other inequalities (Fano, Pinsker, ...)
- many applications

# In Computer Science

- Parallel Repetition

- Communication Complexity



Data Structure,

Streaming algorithms

- Optimization

"Max-Entropy Distribution"

- Crypto, Privacy, ....

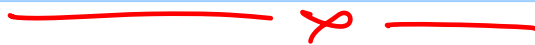
... Punchline

I don't understand any of the above.

Goal of this course:

I should learn.

You will teach me. 😊



So what do I bring to the table

- many authors & course teachers are friends of mine
- I know good papers & can get answers to questions.

# Way the course will work

- Next two lectures:

1 will review information theory basics

- Afterwards ...

- we will (jointly) select papers & present them.

- Each lecture will have a paper + student covering it.

- Student will read paper 1 week before & privately present to me.

We will understand.

Audience will skim paper & come in.

We will discuss / present paper together.

- First paper

"Entropy = Counting" - J.R.

## General Info

Lecturer: MADHU SUDAN

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## Class Info

<http://madhu.seas.harvard.edu/courses/spring2016>

## Do Join Piazza

### To do items

- Sign up for scribing
- Paper + Date of Presentation.
- Questions / Partner?

Use Piazza

Come to OH.